Fast Computation with Efficient Object Data Distribution for Large-Scale Hologram Generation on a Multi-GPU Cluster

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SUMMARY The 3D holographic display has long been expected as a future human interface as it does not require users to wear special devices. However, its heavy computation requirement prevents the realization of such displays. A recent study says that objects and holograms with several giga-pixels should be processed in real time for the realization of high resolution and wide view angle. To this problem, first, we have adapted a conventional FFT algorithm to a GPU cluster environment in order to avoid heavy inter-node communications. Then, we have applied several single-node and multi-node optimization and parallelization techniques. The single-node optimizations include a change of the way of object decomposition, reduction of data transfer between the CPU and GPU, kernel integration, stream processing, and utilization of multiple GPUs within a node. The multi-node optimizations include distribution methods of object data from host node to the other nodes. Experimental results show that intra-node optimizations attain 11.52 times speed-up from the original single node code. Further, multi-node optimizations using 8 nodes, 2 GPUs per node, attain an execution time of 4.28 sec for generating a 1.6 giga-pixel hologram from a 3.2 giga-pixel object. It means a 237.92 times speed-up of the sequential processing by CPU and 41.78 times speed-up of multi-threaded execution on multicore-CPU, using a conventional FFT-based algorithm.

key words: computer generated holography, large-scale CGH, GPU cluster

1. Introduction

The 3D holographic display has long been expected as a future human interface as it does not require users to wear special devices [1]. In display systems, the computer receives 3D object data, computes wave propagation from each object point to each hologram point as well as the interference with a reference laser beam and stores the whole result to a SLM (Spatial Light Modulator). By applying a reference laser illumination to the SLM, the diffraction pattern regenerates the object wave and the user can see the reconstructed 3D object as if it exists at the original position. This process is well-known as computer generated holography (CGH) [2], [3].

One of the largest issues for realizing such display systems is their requirement for a huge number of computations. The computation cost of wave propagation from N object points to N hologram points becomes \(O(N^2)\) [9]. In order to solve this problem, a wide variety of approaches have been taken. To accelerate the calculation of point-to-point wave propagation, a table-lookup method is used to evaluate trigonometric functions [4], [16]. Architectural approaches include the development of special-purpose computers using large-scale FPGA chips [10], [11] and the utilization of GPU and GPU clusters [9], [12]–[15], [18], [25]. In order to reduce the large cost of point-to-point calculation, FFT-based methods are used to improve the order of computation from \(O(N^2)\) to \(O(N \log N)\) [17]. However, for the size of the objects and holograms these projects treat, the processing time for one frame becomes hundreds to even thousands of seconds [12], [26].

On the other hand, the capability of 14.7 giga pixel object display is required for the human eyes’ recognition capability [7]. Further, a 4.5 giga-pixel SLM is needed to realize view angles of 20 degree for 7 inch objects [8]. Thus, the giga-pixel order is required to realize the high resolution and wide view angle needed to match the human interface.

If we treat, for example, 4 giga-pixels, each pixel requires 8B as a complex number representation and a total of 32 GB of memory space is required to store just the object data. We need additional memory space for the output hologram and working space. The global memory size of current GPUs is far behind these requirements [28]. Thus, if we want to utilize the parallel computation capability of GPU clusters as well as the low computation cost of FFT based methods, we need to decompose the input object into multiple sub-objects and apply distributed FFT algorithms to them. In this case, we encounter a further problem that the distributed FFT algorithms require a lot of node-to-node data transfer for realizing butterfly operations [18].

To this problem, we have been working on a distributed algorithm, called a decomposition method, for generating 2D Fourier holograms [5] and 3D Fresnel holograms [6]. In order to estimate the performance of the method, we per-
formed a simulation. In this simulation, as we could not use GPU-clusters, we first obtained computation time on a single multi-GPU node without application-oriented optimizations and inserted the time to the CPU-cluster program by using the nano-sleep library function. The simulation results were used for showing the advantage of our decomposition method over conventional methods. In parallel with the simulation, we have also developed a real GPU cluster using 8 nodes of multi-GPU machines, implemented the decomposition method for 3D Fresnel hologram generation on the cluster and applied application-oriented optimizations [19]–[21]. Thus, the major differences of this research from [6] are that this research utilizes a real multi-GPU cluster and applies various application-oriented, single-node and multi-node optimizations to verify their effectiveness.

Aiming for the ultimate goal of realizing 3D holographic display with high-resolution and wide view angle properties in real time, this research revises the results of [19], [20] and [21] and shows how we resolve the difficulties of large-scale CGH generation on a multi-GPU cluster by adapting the FFT-based algorithm to the clusters’ environment and how to apply application-oriented optimizations under the multicore-CPU and multi-GPU combined heterogeneous architecture.

This paper is organized as follows. Section 2 describes the basic concept of the conventional FFT-based CGH generation and the object decomposition method. Sections 3 and 4 describe the application of optimizations within single node and multiple nodes, respectively. Section 5 shows experimental results using a multi-GPU cluster. Section 6 concludes the paper.

2. Algorithm Adaptation for GPU Cluster Implementation

Figure 1 shows the process of 3D image reconstruction using conventional FFT-based methods. Notice that a 3D object is represented as multiple 2D planes, called object layers, in order to apply 2D-FFT operations. In the figure an input object consists of two layers placed at different positions. Firstly, the wave propagation from these multiple object layers to the hologram plane is computed by applying 2D-FFT operations; then, the interference with the reference laser beam is added to make hologram data, called interference fringes; then, this hologram data is transferred to the SLM; finally, by illuminating the SLM with the reference laser light the diffraction regenerates the object wave from the 3D object. This regenerated wave enables the viewer to see as if the original object exists at the original position.

Equations (1) to (3) below express the wave propagation computation from an object layer \( O(x_1, y_1) \) to the hologram plane \( H(x, y) \).

\[
H(x, y) = QP(x, y) \times FFT_{N \times N}[O(x_1, y_1) \times QP_1(x_1, y_1)]
\]

\[
QP(x, y) = e^{j \frac{kx^2 + y^2}{2z}}
\]

\[
QP_1(x_1, y_1) = e^{j \frac{k(x_1 + y_1)^2}{2z}}
\]

In these expressions, \( H \) represents a hologram plane, \( x \) and \( y \) represent the hologram plane’s horizontal and vertical coordinates, respectively, \( O \) is the two-dimensional object layer, \( x_1 \) and \( y_1 \) are the object’s horizontal and vertical coordinates, \( z \) represents the distance from the object layer to the recording hologram plane, \( N \) is the number of pixels on one dimension of the object, and \( k \) is the value of \( 2\pi \) divided by the wavelength of light (\( \lambda \)).

Figure 2 shows the flow of conventional FFT-based 3D Fresnel hologram computation. The five steps from “Generate \( QP \)” to “Multiply \( QP \)” realize Eq. (1). That is, “Generate \( QP \)” and “Calculate \( O \times QP \)” compute \( O(x_1, y_1) \times QP_1(x_1, y_1) \). The next two steps “2D-FFT” realizes \( FFT_{N \times N} \). The next two steps of “Generate \( QP \)” and “Multiply \( QP \)” realize multiplication by Eq. (2). Notice that the first two steps of “Generate Random num.” and “Randomize” are inserted as an optical convention so that light from all parts of the object can spread over the hologram plane [22]. The step “Add reference beam” implements the interference with the reference light. The “Final addition” accumulates all the results for different layers.

Figure 3 shows the process of hologram computation and 3D image reconstruction using the object decomposition method [5]. In the figure, the input object consists of two 2D object layers, a pigeon and an olive. Each object layer is first decomposed into sub-objects; in the figure, we assume the use of 4 nodes and decompose the object layer into 4 sub-objects; then the three operations of interpolation,
FFT, and shift are applied to each sub-object to produce a hologram, called a sub-CGH; by adding two sub-CGHs for sub-objects of different layers we can obtain the final sub-CGH; and by illuminating this final sub-CGH the original sub-objects are reconstructed at their original positions. **Interpolation** is necessary to keep the original bandwidth and make the object segment reconstructed successfully. A **shift** operation is necessary for reconstruction at the original position of the object segment. Without this shift, all the object segments are reconstructed at the upper-left corner. Theoretically, if we add up all the sub-CGHs, we can obtain the same results with those obtained by the conventional FFT-based method [17]. However, in order to save this addition time, we suppose the display system which switches its input from among the generated, multiple sub-CGHs. This enables us to reconstruct the original 3D object by time sequential reconstruction of sub-objects. We have proved that if the time sequential reconstruction is performed at an appropriate speed, the whole original 3D image is perceived [5, 6].

Notice that the boxes in Fig. 3 indicate that after the original object is decomposed and distributed to parallel processing units, the computation from sub-object to sub-CGH and reconstruction can be performed in a mutually independent manner.

For the theoretical background and optical verification of the decomposition method as well as the preliminary results of its application to the Fourier and Fresnel hologram generation, see [5] and [6], respectively.

Figure 4 shows the basic flow of the decomposition method that generates each sub-CGH from each sub-object. In the figure, the interpolation, 1D-FFT and shift operations are applied in the x- and y-directions, sequentially. This process is repeated for different object layers and the results are accumulated by the final addition. Notice that the functions from “MemcpyHtoD” to “MemcpyDtoH” are executed in the GPUs. Thus, the time-consuming “Add reference beam” should be included in these functions.

At first, we have implemented this flow on a single-node with multiple GPUs as the base program of further optimizations and parallelization. Considering the limited global memory size of each GPU, we have tried to off-load the sub-object to sub-CGH generation, the most heavy computation part, to the GPUs. The CPU controls the global flow and does memory intensive work, i.e. keeping the whole object and the generated sub-CGH.

Figure 5 shows an illustration of the sub-CGH generation process by the loop iterations in the flowchart. At the first iteration, the application of the x and y direction operations to each sub-object produces the same size of partial sub-CGH, called sub-sub-CGH data. A GPU performs this part of the computation K times, where K represents the number of decompositions, and sends the results to the CPU. Then, the CPU adds the sub-sub-CGH data to the predetermined area, as shown in Fig.5(a). In the figure, j is used as an index for counting loop iterations up to the number of decompositions, i.e., K.

By filling up the sub-CGH area in CPU memory by
four times of sub-sub-CGH data, sent from GPU, the sub-CGH is completed. Notice that the number of loop iterations reflects the number of decompositions. The $x$- and $y$-direction operations should be repeated by the number of $x$- and $y$-direction decompositions, respectively. For example, if the input object is decomposed like Fig. 3 the numbers of $x$- and $y$-direction iterations become 2 and 2, respectively, and a total of 4 sub-sub-CGH data generation processes are performed.

This means that the increase of the number of decomposition increases the number of loop iterations and thus the total computation cost. Therefore, from the view point of computation cost, a small number of decompositions is good. On the other hand, from the view point of global memory size, a large number of decompositions is good as it makes the sub-object size small. Thus, in general, we should select the number of decompositions as a tradeoff point of these two aspects.

Comparing with the conventional method, the decomposition method requires less communication as the computation processes from sub-objects to sub-CGHs are mutually data-independent. Thus, once the original object is decomposed into sub-objects and sent to computing nodes, no node-to-node communications are necessary at all. Notice the final reconstructions are also mutually independent. Another important aspect of the decomposition method is that it requires much less GPU memory space, since just the decomposed sub-object and the corresponding partial sub-CGH data are treated in the GPU. This is critical for generating gigapixel hologram using a limited global memory space of GPU.

However, we should be careful about the following two points: first, the decomposition method increases the amount of computation for interpolation and space-shift; second, the sub-CGH generated from the sub-object requires the same size of memory space as the whole CGH.

Thus, the key issues of object decomposition method implementation on GPU clusters are to determine: i) whether the additional computation cost for the interpolation and shift operations is paid back by less communication and by the deletion of the addition of sub-CGHs, and ii) whether the sub-sub-CGH data is generated within a limited global memory space of GPU and appropriately stored back to the host CPU to make sub-CGH used for the reconstruction. Our effort of optimizations and parallelization (Sects. 3 and 4) and the experimental results (Sect. 5) will answer these issues.

3. Single-Node, Multi-GPU Optimizations

This section describes the optimizations and parallelization within a single node with multiple GPUs.

3.1 Changing Distributed Accesses to Contiguous Ones (Optimization 1)

When we send data from the CPU to a GPU, the data transfer is performed by the cudaMemcpy function, which requires the input data to be on a contiguous address space.

At first, we decompose the input object plane into $x$- and $y$-directions like shown in Fig. 6 (a). In this case, we need to rearrange the data so that the selected parts become contiguous in memory before transferring from CPU to GPU.

We have improved the way of decomposition by changing to $y$-direction only decomposition, as shown in Fig. 6(b). This makes the decomposed areas contiguous in memory and we can transfer the data without rearrangement.

Figure 7 shows the flow after this optimization. By comparing this flow with the original flow of Fig. 4, we can understand that, in the original flow, the combination of the $x$- and $y$-direction operations is repeatedly executed. However, in Fig. 5, once the $x$-direction operation is executed, the resulting data in local memory can be utilized repeatedly by the following $y$-direction operations. This not only reduces the number of kernel executions for $x$-direction operations but also enables the passing of temporal computation results from a $x$-direction computation to a $y$-direction one through high-speed local memory.

Notice that in the original way of decomposition the number of decompositions should be $N^2$ where $N = 1, 2, \ldots$. 
However, this change of the way of decomposition alters the number to just $N$ and thus makes it easier to adjust the sub-object size to the limitation of a GPU’s memory size. In our implementation, for the ease of coding, $N$ is selected as $2^n$ where $n = 1, 2, \ldots$. 

3.2 Reduction of Data Transfer Amount between CPU and GPU (Optimization 2)

As the data transfer time between CPU and GPU adds to the computation time, the GPU’s fast computation may be impaired by the slow transfer time. Thus, we should be careful about the contents of the transfer.

Our CGH generation program uses cufft, a CUDA library, for the FFT computation [23]. It requires cuFFTComplex as its input array. The cuFFTComplex consists of a real part and an imaginary part, and each part requires 4B memory space for storing floating point data. Thus, in our original program, both CPU-to-GPU input transfer and GPU-to-CPU output transfer use an 8B cuFFTComplex type.

However, the use of this data type is redundant. A pixel of the input object uses just one byte of the real part. This can be replaced with 8-bit unsigned integers. In a similar manner, the output from GPU only uses a 4B real part and can be replaced with a floating point number.

Figure 8 shows this change schematically. The reduction of the data size shortens the data transfers from CPU to GPU and from GPU to CPU. However, the GPU must perform the additional tasks of extension from the input 8-bit to 8-byte complex type and reduction from 8-byte complex type to 4-byte real type. Experimental results are expected to clarify the total effect.

3.3 Kernel Integration and Utilization of High-Speed Memory (Optimization 3)

Usually, a GPU has thousands of execution cores for arithmetic and logical operations [24]. In order to utilize their computing power effectively, it is critical to smoothly provide them data to be processed. For this purpose, a GPU has a hierarchically structured memory system of register file, local memory, cache memory (usually multi-level), and global memory. The input data to a GPU is first sent to the global memory. When a kernel starts its execution, it reads data from the global memory to a higher level memory and tries to keep the data at the higher level until the end of the execution. At the end of the execution, the results should be saved to global memory to pass them to the succeeding kernel. Then the succeeding kernel reads the data again into higher level memory to perform its operation. The problem here is the large overhead of this write-back and read-out cost of global memory, the slowest memory in the hierarchy. We cannot overlook that these write-back and read-out operations also lose the locality of references.

In order to solve this problem, we have tried to integrate kernel functions into one kernel function as much as possible. Figure 9 shows the integrations of the three kernels from “Randomize” to “Calculate D x QP” and the five kernels from “Shift” to “Add reference beam”. These kernels are called separately from the CPU in Fig. 7. Before the integration, each kernel reads all the pixels from global memory, applies each computation and stores the results back to global memory. After the integration, the three kernel functions are combined into one kernel function. This kernel function reads a part of the pixels from global memory, applies the three operations to them, i.e. Shift, Transpose and Multiply QP, and stores the partial results back to global memory, repeatedly, until all the pixels are processed. By this integration, we can utilize higher-level memories, such as registers and shared memory, for inter-kernel data passing. The integration also reduces the number of kernel function calls and, thus, reduces the cost for GPU control.

3.4 Stream Processing of GPU Calculations and Transfer of the Result (Optimization 4)

As shown in Fig. 5, the sub-CGH, generated from the sub-object, becomes the same size as the final CGH. Thus, we need to send the partial results, i.e., sub-sub-CGH data, back to the CPU repeatedly to save the limited memory space of the GPU. Fortunately, the generation process can be divided...
into $K$ mutually data-independent operations where $K$ represents the number of decompositions. Thus, we repeat the process $K$ times.

We have utilized stream processing to treat these computation and communication pairs efficiently, as shown in Fig. 10. That is, first, the pair of the partial result calculations in the GPU, represented as A and B, are executed. Then, the results are transferred from the GPU to the CPU by sub-CGH data transfer operation, represented as C. Finally, the final addition operation, D, is executed in CPU.

The pipelining of these operations should take into account the following conditions. The first pair of calculations should be serial as B inputs the output of A. The transfer from GPU to CPU should wait for the completion of B in the GPU. The calculation at A should wait for the completion of B of the preceding pipeline as they share the same buffer memory.

3.5 Utilization of Multiple GPUs under the Control of a Multicore CPU (Optimization 5)

If multiple GPUs are available in a single node, we may utilize their parallel processing capability. In this case, it is desirable to assign them mutually data-independent processing so that they can execute in parallel.

In the CGH generation, if we assign sub-objects of different layers to the multiple GPUs, the generation process becomes mutually independent. In order to control these parallel operations in multiple GPUs, we use multithreading in the CPU. One thread controls one GPU, by specifying the GPU using cudaSetDevice [23].

One important issue for utilizing the multiple GPUs is how the CPU receives the results of the computation from them and, if necessary, does some reduction operations on the results. In our decomposition model, the reduction operation is the addition of sub-CGHs generated from different object layers. As shown in Fig. 11. We considered three methods to treat this issue: (i) providing disjoint areas for different GPUs, (ii) using “lock” to prevent the other threads from modifying the same area, and (iii) using a synchronization mechanism to avoid allowing the two threads to handle the shared data area simultaneously. We performed preliminary evaluation of these methods and found that (iii) attains the best performance [19]. Thus, we use this method for our experiments, as described in Sect. 5.

4. Multi-Node Optimizations

The major principle of the decomposition method is to save communication cost by avoiding the frequent data transfers required for 2D-FFT operations. Following this principle, we assign sub-objects with the same x- and y-coordinate values but with different layers to a single node. This assignment avoids unnecessary data transfers during the computation.

One inevitable data transfer operation occurs when the host node distributes sub-objects of different layers to the other nodes. We must consider that the input object is given in a compressed image file, such as JPEG and PNG. The object in the compressed format can not be decomposed and used for further computation and we need to decompress it before decomposition. Thus, there are two possibilities.

The first method is to decompress the input object to bitmap, i.e., one pixel to 8-bit unsigned integer, on the host node, decompose the decompressed object into sub-objects, and transfer each sub-object to an appropriate node. In this method, the transfer amount, i.e., the sub-object size, is determined statically. Thus, the size should be smaller than the size of the original compressed object. Figure 12 shows how the input object is decompressed and decomposed into sub-objects in the host node, Node 0, and sent to the other nodes for sub-CGH generation. The decompressed and decomposed sub-objects are sent by MPI send and recv pair operations, sequentially [27].

The second method is that the host node first distributes the compressed whole object to all the other nodes, as shown in Fig. 13. Then, each node decompresses the input object and extracts the sub-object to be processed in the node. In this method, the data size of the compressed object is dynamically determined. Thus, the size should be distributed first from the host to the other nodes so that they can allocate necessary memory space for storing input compressed object. The MPI broadcast is used to distribute the compressed whole object. Each node needs to perform decompression separately. For the details of the sending procedure by using MPI, see [19].

In the following section, we will show the performance
results of these two methods.

5. Experimental Results

We have performed experiments in order to verify the effectiveness of the optimizations. The execution environment is summarized in Table 1. Simply speaking, it is 1000BASE-T connected 8 node GPU cluster with each node containing two GPUs. Many similar projects use a high-speed network, such as Infiniband, to treat this kind of large-scale bandwidth intensive task without locality [18], [25]. Unlike them, we use a low cost gigabit Ethernet. We use C to describe CPU code, Pthreads to describe multithreading for optimization 5, MPI Ver. 1.10.3 to describe the internode communication, and CUDA version 8.0 to describe the GPU code.

The input object consists of a pair of PNG files (40K for each) and the total data size is 141.3MB. They are decompressed into two bitmap files (1.6GB for each) by using an OpenCV function. Thus, the total size of the input object is 3.2 giga-pixels. The hologram size is 1.6 giga-pixels.

Following the discussion in Sect. 2, the number of decompositions is determined to be 8 as a trade-off between the decomposed object size and computation cost. Our first experiment is to apply the optimizations 1 through 5 to clarify their effect in a single node. Throughout the experiment the number of decompositions is fixed at not 8, as determined above, but 16. The reason is that the number should be $2^n$ before optimization 1, as noted in 3.1, and we must use the same value consistently during single-node optimizations. Thus, the host CPU sends a total of 32 (= number of decompositions) $\times$ (number of layers) = 16 $\times$ 2 sub-objects to the GPUs to process all of them.

Figure 14 shows the execution time results where the time of each bar is decomposed into the times for decomposition, memory copy (memcpy), kernel, and the final addition. The kernel processes 32 sub-objects and the kernel time includes the kernel-call time.

The bar A represents the starting point of our optimization. Conventional optimizations, such as an appropriate definition of grid and blocks, coalescing access of global memory, and efficient use of local memory, have already been applied.
The bar $B$ shows the effect of optimization 1. The change in decomposition method eliminates the decomposition time, shown as a very small white box at the top of bar $A$. It also reduces the kernel time as the results of the $x_1$ direction processing is repetitively utilized as described in 3.1. Notice that the effect of results output as contiguous memory space also shortens the final addition time of CPU.

The bar $C$ shows the effect of optimization 2. The results show that the increased kernel time reduces memory copy time. Changing from a complex data type to a real one also reduces the time for the final addition on the CPU.

The bar $D$ shows the effect of optimization 3. The integration of kernels naturally reduces the kernel time.

As to the bars $E$ and $F$, we cannot simply decompose the execution time as the optimizations realize parallel operations between computation and communication (optimization 4) and between multiple GPUs (optimization 5).

The bar $E$ shows the effect of stream processing. The measured time shows improvement, but not as much as we initially expected. The reason is twofold: first, in order to save the memory space of GPU, we did not utilize double buffering between the four operations, and second, the execution times of four pipe-stages are imbalanced (the execution time ratio of the pipe-stages $A$, $B$, $C$ and $D$ in Fig. 10 was $1:1:2:4$) and the fourth stage of addition by CPU becomes the bottleneck.

The bar $F$ represents the effect of optimization 5. Using two GPUs nearly halves the time for $E$.

Eventually, $F$ attains a 11.52 times speed-up from $A$.

Figure 15 shows the execution times of the CPU and the GPU cluster, while changing the number of cluster nodes.

The left-most bar represents the sequential processing time by the CPU, following the flow of Fig. 2, i.e. a conventional FFT-based algorithm. The CPU uses just the CPU part of Table 1. The sequential code is written in C and the 2D-FFT is realized by the FFTW 3.3.2 library. We obtain this time as the baseline performance. Thus, we intentionally do not utilize CPU parallelism, such as multicore multi-threading and SIMD extensions. The results show 1018.32 seconds of processing time.

For the experiment using the 8 node GPU cluster, we fixed the number of decomposition at 8. This is based on the discussion in Sect. 2.

For the next bar of 1 node we use one node of the GPU cluster defined in Table 1. From 2 nodes to 8 nodes, we applied the multi-node optimizations described in Sect. 4. The bars for (a) and (b) show the execution time for the decompress-and-unicast method and multicast-and-decompress one, respectively. Each bar is decomposed into CGH generation time and transfer one.

From the results we understand that broadcasting compressed object is much faster than sequential unicasting of decompressed sub-objects. The reason is as follows: the broadcasting transfers the same data to all nodes and thus the total transfer amount of the method (b), i.e., compressed object, became 141.3 MB (i.e., size of 2 PNG files); on the other hand, that of method (a), i.e., decompressed and decomposed data for 8 nodes, uses multiple unicasting and its total transfer amount became 2.8 GB ($= (size of decompressed and decomposed object) \times (number of nodes - 1) = 400MB \times 7$). Thus, the method (a) is even slower than the execution time by a single node.

From the viewpoint of scalability, the increase of nodes decreases the generation time but at the same time increases the transfer time. Even in method (b), the transfer time slightly increases. These results indicate that the bottleneck of the current system exists in the transfer time.

The rightmost bar indicates that 8 nodes cluster attains the execution time of 4.28 sec, which means 237.92 times speed-up of the sequential processing on CPU.

Further, for fair comparison with a CPU without GPUs, we have performed an experiment of multi-threading on a multicore-CPU using OpenMP. The program has been compiled by gcc ver.4.8.5(CentOS 7's default compiler) with optimization option -O3. As the multi-node implementation is slower than the single node one due to its heavy node-to-node communication costs and low communication speed of 1000BASE-T network, we show the single CPU results in Fig. 16.
The results show that the case for 12 threads attains the shortest execution time of 178.81 sec. This means that our results attain 41.78 times better performance than that for a multicore-CPU.

6. Conclusion

We have described our research results for overcoming the difficulty of large-scale CGH generation on a multi-GPU cluster. Our efforts include algorithm adaptation for a multi-GPU cluster and both intra- and inter-node optimizations of the code for multicore CPU and multi-GPU combined heterogeneous node architecture.

The intra-node optimizations attain an 11.52 times speed-up from the original single node code. The extreme results, using our 8 nodes 2-GPU architecture, show a 4.28 sec execution time for generating a 1.6 giga-pixel hologram. This is 237.92 times faster than the multithreaded implementation.

Thus, the major contribution of this paper is to show that we can generate a giga-pixel CGH from a giga-pixel object in a few seconds under the constraints of limited memory size of GPUs. It is enabled by adapting an FFT-based algorithm to GPU cluster environment and by applying application-oriented optimization and parallelization techniques.

Our future plan is to further reduce the extreme processing time of 4.28 sec to realize our ultimate goal of real-time 3D display with high-resolution and wide view angle. For this objective, we need to analyze the current results and seek for possibilities of performance improvement by using higher performance processors and network as well as their code optimization and parallelization.

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References


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